## 2014 Putnam Problem 3B

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**Proposition 1.** Suppose that A is an  $m \times n$  matrix with rational entries such that the matrix of absolute values of A,  $\widetilde{A}$ , contains at least m + n distinct primes. Prove that A has rank at least two.

*Proof.* Note that it suffices to show that not all rows of A are scalar multiples of the first row of A.

Throughout this problem, we let S denote a set of m+n distinct primes in  $\widetilde{A}$ . First, we prove a lemma:

**Lemma 2.** Let G be a graph on n vertices with n edges. Then, G contains a loop.

*Proof.* Proceed by induction. The case n=1 is trivial. Now, consider a graph G with n+1 vertices and n edges. By the pigeonhole principle, G must contain a vertex v of index at most one. Excise v and its edges from G. If v has index zero, remove v and any edge. What remains is a subgraph with n vertices and n edges. By our inductive hypothesis, this subgraph contains a loop.  $\square$ 

To proceed, we will first solve the case when A is  $n \times n$ , then show that this argument effectively handles the general  $m \times n$  case.

Let A be an  $n \times n$  matrix satisfying the conditions of the proposition. Define a graph G(A) on A as follows. Draw a vertex  $v_i$  corresponding to each column  $c_i$  for  $1 \le i \le n$ . Now, define a function on the rows of A as follows. Let

$$p: \{ \text{rows of } A \} \longrightarrow \mathbf{Z}_{\geq 0}$$
  
 $p(r_i) = \# \{ \text{elements of } S \text{ in row } i \}$ 

Begin in row 1. Consider the  $C_1$  of columns containing an element in S is row one. Draw a connected tree between the vertices corresponding to columns in  $C_1$ . Note that this tree has  $p(r_1) - 1$  edges. Do the same for row j, for  $2 \le j \le n$ , adding the new edges onto edges we have already drawn. Note that each row  $r_i$  adds  $p(r_i) - 1$  edges.

Decompose S as follows. For  $i=1,\ldots,n$ , let  $S_i$  be the primes in S contained in row i of A. Note  $S_i \cap S_j = \emptyset$  for  $i \neq j$  and  $S = \bigcup_{i=1}^n S_i$ .

**Lemma 3.** G(A) contains a loop.

*Proof.* We will now count the number of edges in G(A), which we will call E(G(A)).

$$E(G(A)) = \sum_{i=1}^{n} p(r_i) - 1$$

$$= \sum_{i=1}^{n} p(r_i) - \sum_{i=1}^{n} 1$$

$$= 2n - n$$

$$= n$$

It follows from lemma 2 that G(A) contains a loop.

Now, suppose that all columns of A are scalar multiples of the first colum (that A has rank at most one). Then, when two vertices  $v_i, v_j$  are connected by an edge in G(A),  $c_i = \lambda c_j$ , where  $\lambda$  can be expressed as a ratio of elements of S that live in the same row of A.

Choose a vertex  $v_i$  contained in a loop of G. Then, there are scalars  $\lambda_0, \lambda_1, \ldots, \lambda_k$  such that

$$c_i = \lambda_1 \dots \lambda_k c_j$$
$$c_i = \lambda_0 c_j$$

where  $|\lambda_i|$  is a ratio of distinct primes in S.

We will derive a contradiction from the equation

$$|\lambda_0| = |\lambda_1 \dots \lambda_k| \tag{1}$$

By performing an elementary row operation, we may assume without loss of generality that  $\lambda_0$  is the ratio of two primes in  $S_1$ , say  $p_0/p_1$ .

Our graph was constructed so that if  $\lambda_{i_1}, \ldots, \lambda_{i_r}$  are ratios of primes in  $S_j$ , then  $\lambda_{i_1}, \ldots, \lambda_{i_r} \neq 1$ . In particular, this product, in lowest terms will contain at least one primes in  $S_j$  in both its numerator and denominator.

Also, by construction, product  $\lambda_1 \dots \lambda_k$  must contain a  $\lambda_k$  such that  $\lambda_k$  is a ratio of primes in  $S_i$  for i > 1. This is because row 1 induced no loops in the graph. By the previous observation, one such ratio  $p_2/p_3$  remains when  $\lambda_1 \dots \lambda_k$  is expressed in lowest terms. Clearing the denominator, we have

$$p_3p_0p_1^*\dots p_s^* = p_1p_2p_1'\dots p_s'$$

where the other primes are in S. This contradicts the fundamental theorem of arithmetic.

This proves the  $n \times n$  case. Note that in an  $m \times n$  matrix for m > n. Then, there must be a row that contains 1 or fewer elements of S. Perform a row operation to move this row to the bottom and consider the upper  $m-1 \times n$  submatrix. Repeat until we have an  $n \times n$  matrix containing 2n entries and apply the preivous proof.